

Scientific Hub of Applied Research in Engineering & Information Technology

Received: 26.05.2022 Revised: 09.06.2022 Accepted: 15.06.2022



Research Article Data Encryption using Graph Domination

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In this modern world, every industry and organization has tough competition to sustain. So, everyone needs to develop their own ideas to keep their organization exists in good condition. The new technology, innovation and updates are very confidential until it will be declared officially. It is mandatory to maintain the data and communicate the information safe from hackers and other external sources. Encryption and decryption help us to prevent and send the details in safe mode. Using graph theory in cryptography is developing as a promising field. Domination theory in the encryption process is rarely used. A new data encryption method using graph function and control number is discussed in this paper. A sample output using MATLAB program also provided using the same procedure.

Keywords: Encryption, Decryption, Domination number, Graph theory.

1. Introduction

Encryption is the method of interpreting plaintext into something that seems to be meaningless, which may be called ciphertext. "The method of converting ciphertext into plaintext is called decryption. Symmetric encryption can be used to encrypt a large amount of data. Using symmetric keys plays an important role in both the process of encryption and decryption. To decrypt a specific portion of ciphertext, the symmetric key that was used to encrypt the message must be used. The main aim of every encryption algorithm is to generate a new method and it will make complicated to decrypt the ciphertext as much as possible without using the encryption key" [3]. Graphic theory plays a major role in encryption, using graph theory in cryptography has gained a lot of attention these days.

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An encryption algorithm using graph theory properties has provided by Etaiwi. Define encryption algorithm and secure message encryption using graphical theory features [8]. In 2013, M. Yamuna et al. provide encryption and decryption of any chemical formula using stable domination separation graphs [10]. Also, they developed a new way to encrypt any tree using a new periodic binary table [9]. A method of data encryption using a circuit matrix was discussed in [5].

2. Preliminary note

The basic explanations and results related to the graph theory required for the proposed encryption system are described in this section.

2.1. Graph

Graph G = (V, E) is a random pair of vertical V and set ed edges E, which is a subset of the vertex set V feature, i.e. the edge is a two-dimensional event and the relation is represented as a pair of random tweets with respect to the corresponding edge.

2.2. Weighted Graph

Graph G is considered to be a weighted graph when a numerical value is assigned to the entire boundary of G [6]. An example of a weighted graph is given in Fig 1.

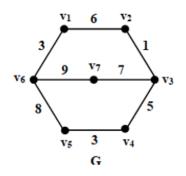


Fig.1. G is a Weighted graph

2.3. Controlling Set

In Graphic G = (V, E), a set of texts D in G is said to dominate the set if each vertex in V - D is close to at least one vertex in D. If D has a small number of elements of any G-ruling set, then D is called the subrule set of G. The amount of elements in the lower Gruling set is called the G-ruling number and is defined as γ (G). The governing set with a small number of G elements is called γ - set [2].

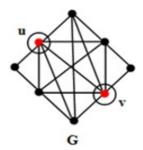


Fig.2. { u, v } is a γ - set for a graph G and γ (G) = 2

Rule D is said to be the independent rule set if there are no two situations in D nearby. Vertex v is said to be a low (critical) vertex if γ (G - v) < γ (G).

3. Proposed approach

Data encryption using graph function and rule number is discussed in this section.

3.1. Edge Contraction

Identifying direct pairs (u v) from G is defined as $G \cdot uv$. Graph $G \cdot uv$ produced by G by subtracting adjacent

vertices (u, v) and attaching a new vertex, say uv, where uv is adjacent to all G vertices - $\{u, v\}$, adjacent or u-u. or v in G [1].

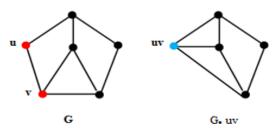


Fig.3. Example of border compaction

Graphic functions are widely used in various research areas. Different kinds of graph operations are available in graph theory, edge contraction is used for encrypting a given message in this paper. Consider the G graph, while the G and G control numbers •UV may or may not decrease when using edge access. If γ (G •uv) = γ (G), then the graph G is called the domination dot stable (DDS) graph [11]. We have graphs, where some edges satisfies the property and no edge satisfies the property. In this paper, the set of edges which satisfy the condition γ (G • uv) = γ (G) as a key for encryption.

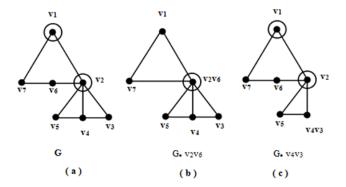


Fig.4. (a) is a graph with γ (G) = 2,

(b) is an example of γ (G• uv) $\neq \gamma$ (G),

(c) is an example of γ (G• uv) = γ (G) = 2

We use the following results to assess the conditions (2)

$$\gamma$$
 (**G**• uv) < γ (**G**) and γ (**G**• uv) = γ (**G**).

Results

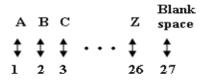
1."Let a, $b \in V (G)$ for a graph G. Then $\gamma (G \bullet ab) < \gamma$ (G) if and only if either there exists a minimum dominating S of G such that a, $b \in S$ or at least one of a or b is critical in G [1].

2. A graph G is DDS if and only if every γ – set of G is an independent dominating set [11]".

3.2. Encryption Chart

To encrypt a message, we can use any standard encryption chart. Sample chart is given in Table 1.

Table 1 Encryption Chart.



We can increase the number of characters invariably depending on the requirements of the message to encrypt.

3.3. Encryption Algorithm

Consider a message that to be encrypted as M. Let us assume the length of M is equal to p.

For example, let M: GRAPHS be the given message to encrypt. Length of M = 6.

Step 1: In the given message M, convert the characters into numerical values using Table 1. Let w_1 , w_2 , w_3 , ..., w_p .

The required weights of the given string M are 7 18 1 16 8 19.

Step 2: Consider a graph G with exactly k - edges satisfying the condition γ (G• uv) = γ (G).

For our example ($v_1 v_2$), ($v_2 v_6$), ($v_3 v_4$), ($v_4 v_5$), ($v_5 v_3$), ($v_6 v_1$) are satisfying the above condition. A graph with exactly six edges satisfy the condition γ (G• uv) = γ (G) as seen in Figure 5 (a).

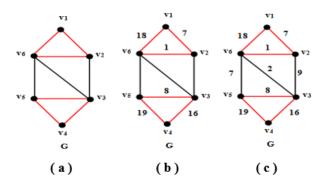


Fig.5. Encryption Algorithm

Step 3: Let X be the set of p edges that satisfy the condition γ (G• uv) = γ (G). Arrange the vertex pairs in the set X in the following order. Let vi be a vertex belongs to the set X, where i is the smallest value in the

set. Collect all possible vertex pairs with vi, say ($v_i v_j$), where i < j and j must be in increasing order. Choose the next smallest value and repeat the same process, stop this process until p edges are selected.

Arrange the vertex pair as discussed above and assign weights { w_1 , w_2 , w_3 , ..., w_k } edges { e_1 , e_2 , ..., e_k }.

For our example { $e_1, e_2, ..., e_6$ } = { ($v_1 v_2$), ($v_1 v_6$), ($v_2 v_6$), ($v_3 v_4$), ($v_3 v_5$), ($v_4 v_5$) } and { $w_1, w_2, ..., w_6$ } = { 7, 18, 1, 16, 8, 19 }.

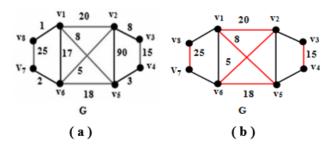
Refer Figure 5 (b) for a graph with weights assigned to the edges selected in step 3.

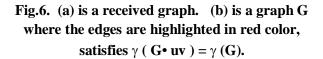
Step 4: Assign random weights to the remaining edges of G.

Step 5: Send the graph G to the receiver as seen in Figure 5 (c).

3.4. Decryption Algorithm

We can decrypt the graph by reversing the procedure into its corresponding data. Consider a received graph G. For example, consider the graph as shown in Fig 6.





Step 1: Choose the edges which satisfied the condition γ (G• uv) = γ (G).

For our example, γ (G) = 3 and the edges ($v_1 v_2$), ($v_3 v_4$), ($v_5 v_6$), ($v_7 v_8$), ($v_1 v_5$), ($v_2 v_6$) which satisfies the condition γ (G• uv) = γ (G) = 3, as seen in Figure 6 (b).

Step 2: List the highlighted edges in increasing order.

For our example, ($v_1 v_2$), ($v_1 v_5$), ($v_2 v_6$), ($v_3 v_4$), ($v_5 v_6$), ($v_7 v_8$).

Step 3: Convert the weights { 20, 8, 5, 15, 18, 25 } into its corresponding character using Table 1.

The required decrypted message is THEORY.

4. MATLAB simulation

The sample output of encryption and decryption using MATLAB program is discussed in Figure 7 and 8.

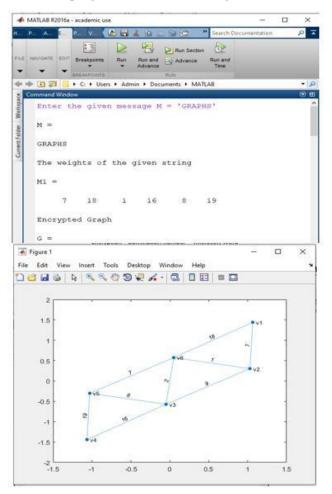


Fig.7. The message and encrypted graph

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Fig.8. The decrypted message

5. Conclusion

Graphs play an important mathematical tool in many real - world problems. Graph theory properties will strengthen the safe transmission. In addition, we have used graph operation and domination property for encryption. The text may be of any length. We can create a graph with any number of vertices and edges to the structure discussed. satisfy Nowadays communicating the details to others is unavoidable. Message transmission using this method is more secure since all the edges and vertices are not used as a key for encryption. Without knowing the domination property, it is very tough to decrypt the given message. A MATLAB code also provided for the same. We conclude that the proposed method will meet the expectations on the safe transmission of any data.

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31